How to Specify it!

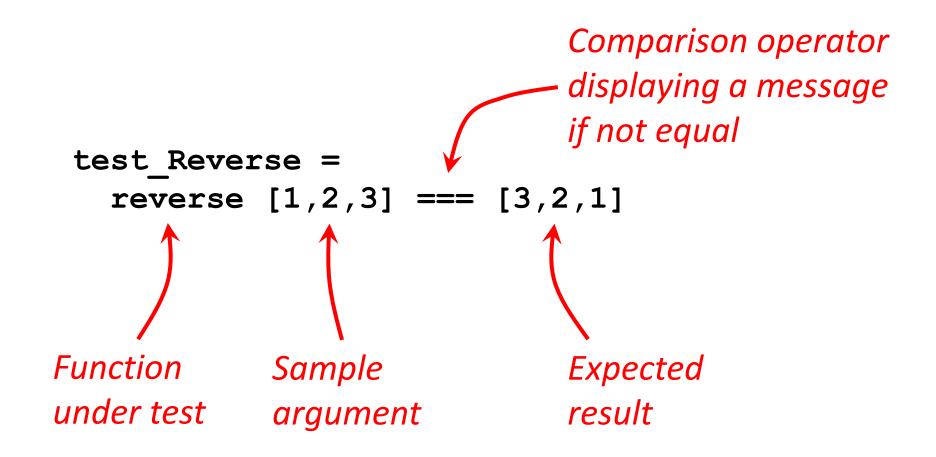
A guide to writing properties of pure functions.

John Hughes





Imagine testing reverse...



Imagine testing reverse...

```
with QuickCheck
                            Tell QuickCheck to
                            generate random
                            lists of integers
 prop Reverse :: [Int] -> Property
 prop Reverse xs =
   reverse xs === ???
     A random
                             But what do we
                             put here?
     argument
```

Imagine testing **reverse**... with QuickCheck

```
prop Reverse :: [Int] -> Property
prop Reverse xs =
  reverse xs === predictReverse xs
This is not easier
                                 We're likely to
                                 make the same
to write than
                                 mistakes
reverse!
```

Replicating the code in the tests...



Expensive!

Low value!

What can we do instead?

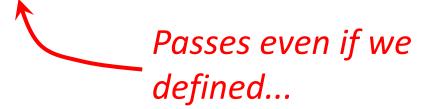
```
prop_Reverse :: [Int] -> Property
prop_Reverse xs =
  reverse (reverse xs) === xs
```

Check a **property** of the return value instead

*Reverse> quickCheck prop_Reverse
+++ OK, passed 100 tests.



reverse xs = xs



*Reverse> quickCheck test_Reverse
*** Failed! Falsified (after 1 test):
[1,2,3] /= [3,2,1]

```
prop_Wrong :: [Int] -> Property
prop_Wrong xs = reverse xs === xs
```

```
*Reverse> quickCheck prop_Wrong

*** Failed! Falsified (after 3 tests and 3 shrinks):

[0,1] 

[1,0] /= [0,1] 

Counterexample:

Almost always [0,1],

sometimes [1,0]
```

Shrinking

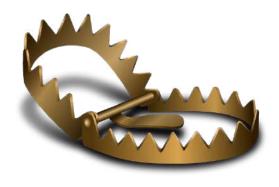
- Discards unnecessary list elements (we need at least two)
- Replaces integers by smaller integers (we need distinct integers, {0,1} are the two smallest)

Property Based Testing

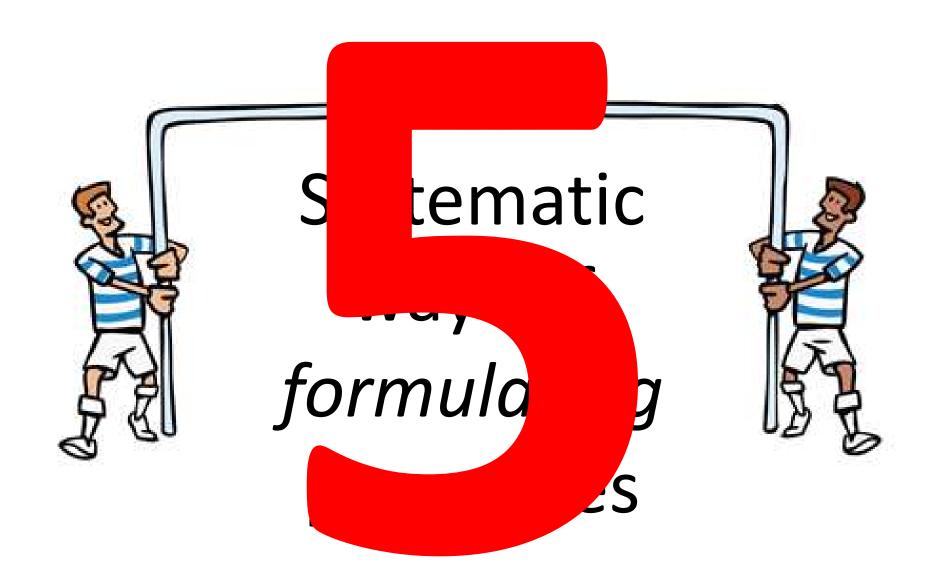


Random generation of *lots* of test cases

Shrinking results in *minimal* counterexamples—easy to debug



 Replicating code in the tests is tempting, but expensive, and low value



Example



```
data BST k v = Leaf
                | Branch (BST k v) k v (BST k v)
  deriving (Eq, Show, Generic)
-- the operations under test
find :: Ord k \Rightarrow k -> BST k v \rightarrow Maybe v
nil ::
                                                    BST k v
insert :: Ord k \Rightarrow k \Rightarrow v \Rightarrow BST k v \Rightarrow BST k v
delete :: Ord k \Rightarrow k -> BST k v \rightarrow BST k v
union :: Ord k \Rightarrow BST k v \rightarrow BST k v \rightarrow BST k v
-- auxiliary operations
toList :: BST k v \rightarrow [(k, v)]
keys :: BST k v \rightarrow [k]
```

Generator and shrinker

```
instance (Ord k, Arbitrary k, Arbitrary v) =>
    Arbitrary (BST k v) where
                                 Generate by
  -- generator
                                 inserting random
  arbitrary =
                                 key-value pairs
    do kvs <- arbitrary
       return $
         foldr (uncurry insert) nil (kvs :: [(k,v)])
  -- shrinker
  shrink = genericShrink
                            Shrink using a
                            generic QuickCheck
                            mechanism
```

1

Is there an invariant?

```
valid :: Ord k => BST k v -> Bool

valid Leaf = True

valid (Branch l k v r) =
 valid l && valid r &&
 all (<k) (keys l) && all (>k) (keys r)
```

Invariant properties

```
prop NilValid = valid (nil :: Tree)
prop InsertValid :: Key -> Val -> Tree -> Bool
prop InsertValid k v t = valid (insert k v t)
prop DeleteValid :: Key -> Tree -> Bool
prop DeleteValid k t = valid (delete k t)
prop UnionValid :: Tree -> Tree -> Bool
prop UnionValid t t' = valid (union t t')
             type Key = Int
             type Val = Int
             type Tree = BST Key Val
```



```
=== prop InsertValid from BSTSpec.hs:19 ===
*** Failed! Falsified (after 6 tests and 8 shrinks):
0
Branch Leaf 0 0 Leaf
=== prop DeleteValid from BSTSpec.hs:22 ===
*** Failed! Falsified (after 8 tests and 7 shrinks):
0
Branch Leaf 1 0 (Branch Leaf 0 0 Leaf)
=== prop UnionValid from BSTSpec.hs:25 ===
*** Failed! Falsified (after 7 tests and 9 shrinks):
Branch Leaf 0 0 (Branch Leaf 0 0 Leaf)
Leaf
```



```
=== prop InsertValid from BSTSpec.hs:19 ===
*** Failed! Falsified (after 6 tests and 8 shrinks):
0
Branch Leaf 0 0 Leaf
=== prop DeleteValid from BSTSpec.hs:22 ===
*** Failed! Falsified (after 8 tests and 7 shrinks):
0
Branch Leaf (1) (Branch Leaf 0)
=== prop UnionValid from BSTSpec.hs:25 ===
*** Failed! Falsified (after 7 tests and 9 shrinks):
Branch Leaf (0 ) (Branch Leaf (0 ) Leaf)
Leaf
```

```
prop_ArbitraryValid t = valid t
```

```
prop_ShrinkValid t =
  all valid (shrink t)
```

Branch Leaf 0 0 (Branch Leaf 0 1 Leaf)
→ Branch Leaf 0 0 (Branch Leaf 0 0 Leaf)

```
prop_ArbitraryValid t = valid t
```

```
prop_ShrinkValid t =
  all valid (shrink t)
```

Branch Leaf 0 0 (Branch Leaf 0 1 Leaf)

→ Branch Leaf 0 0 (Branch Leaf 0 0 Leaf)

```
prop_ArbitraryValid t = valid t
```

```
prop_ShrinkValid t =
  all valid (shrink t)
```

Branch Leaf 0 0 (Branch Leaf 0 1 Leaf)

→ Branch Leaf 0 0 (Branch Leaf 0 0 Leaf)

```
prop ShrinkValid t =
  valid t ==>
    all valid (shrink t)
Branch Leaf 0 0 (Branch Leaf 1 0 Leaf)
→ Branch Leaf 0 0 (Branch Leaf 0 0 Leaf)
```

prop ArbitraryValid t = valid t

Fixing shrinking

```
shrink = genericShrink.
```

What is the *postcondition*?

"After calling insert, we should be able to find the key inserted, and any other keys present beforehand"

What is the postcondition of **find**?

"After calling find,

- —if the key is present in the tree, the result is **Just value**
- —if the key is not present, the result is **Nothing**"

How can we tell this?

By construction!

```
Don't test for Can every tree containing presence—ensure by k be expressed in this construction form?

prop_FindPostPresent k v t = find k (insert k v t) === Just v
```

```
prop_FindPostAbsent k t =
  find k (delete k t) === Nothing
```

3

Metamorphic properties

"How does changing the *input* of insert change its result?" $O(n^2)$ ideas insert k' insert k v insert k insert k' v'

```
Modified call
prop InsertInsert (k,v) (k',v') t =
  insert k v (insert k' v' t)
  insert k' v' (insert k v t)
                            Original call
            Relationship
```

```
prop InsertInsert (k,v) (k',v') t =
  insert k v (insert k' v' t)
  insert k' v' (insert k v t)
                                          TEST
                — Is this really true?
                                          IT!!!
=== prop InsertInsert from BSTSpec.hs:78 ===
*** Failed! Falsified (after 2 tests and 5 shrinks):
(0,0)
(0,1)
                       Last insertion wins
Leaf
Branch Leaf 0 0 Leaf /= Branch Leaf 0 1 Leaf
```

```
prop InsertInsert (k,v) (k',v') t =
  insert k v (insert k' v' t)
  if k==k' then insert k v t else
  insert k' v' (insert k v t)
=== prop InsertInsert from BSTSpec.hs:78 ===
*** Failed! Falsified (after 2 tests):
(1,0)
(0,0)
Leaf
Branch Leaf 0 0 (Branch Leaf 1 0 Leaf) /=
Branch (Branch Leaf 0 0 Leaf) 1 0 Leaf
```

Equivalence for trees

```
t1 =~= t2 =
  toList t1 === toList t2

prop_InsertInsert (k,v) (k',v') t =
  insert k v (insert k' v' t)
  =~=
  if k==k' then insert k v t else
  insert k' v' (insert k v t)
```





Inductive properties

```
prop_UnionNil t =
  union nil t === t
```

```
prop_UnionInsert t t' (k,v) =
   union (insert k v t) t'

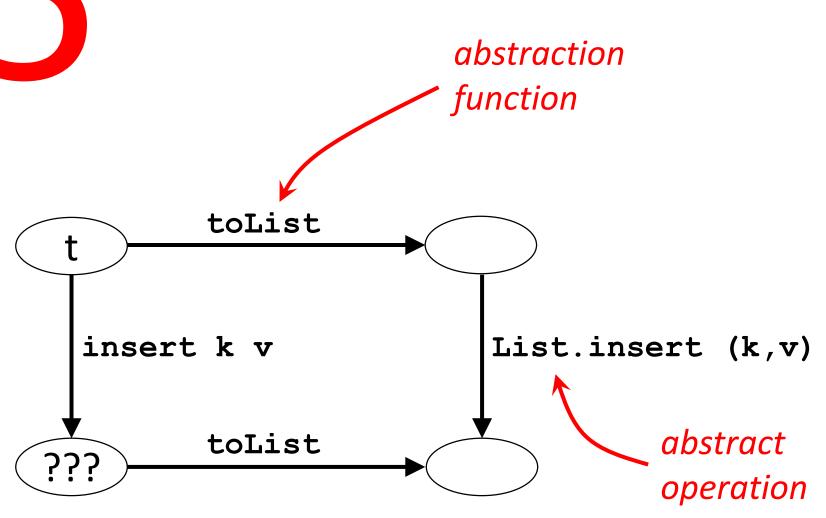
=~=
  insert k v (union t t')
```

union must be correct, by induction on the first argument

Completeness: can every tree be constructed just using insert?

```
insertions Leaf = []
insertions (Branch l k v r) =
  (k,v):insertions l++insertions r
prop InsertComplete t =
                  — Structurally equal to!
  fold1 (flip $ uncurry insert) nil (insertions t)
prop InsertCompleteForDelete k t =
 prop InsertComplete (delete k t)
prop InsertCompleteForUnion t t' =
 prop InsertComplete (union t t')
```

Model-based properties



```
prop InsertModel (k,v) t =
  toList (insert k v t)
  L.insert (k,v) (toList t)
*BSTSpec> quickCheck prop InsertModel
*** Failed! Falsified (after 13 tests and 7 shrinks):
(1,0)
Branch Leaf 1 0 Leaf
[(1,0)] /= [(1,0),(1,0)]
                          · duplicated key
```

```
prop_InsertModel (k,v) t =
  toList (insert k v t)
  ===
  L.insert (k,v) (deleteKey k $ toList t)
```

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Proof of Correctness of Data Representations

C. A. R. Hoare

Received February 16, 1972

Summary. A powerful method of simplifying the proofs of program correctness is suggested; and some new light is shed on the problem of functions with side-effects.

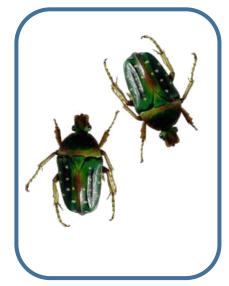
1. Introduction

In the development of programs by stepwise refinement [1-4], the programmer is encouraged to postpone the decision on the representation of his data until after he has designed his algorithm, and has expressed it as an "abstract" program operating on "abstract" data. He then chooses for the abstract data some convenient and efficient concrete representation in the store of a computer; and finally programs the primitive operations required by his abstract program in terms of this concrete representation. This paper suggests an automatic method of accomplishing the transition between an abstract and a concrete program, and also a method of proving its correctness; that is, of proving that the concrete representation exhibits all the properties expected of it by the "abstract" pro-

Type of property	Number of properties
Invariant	4
Postcondition	5
Metamorphic	16
Model-based	5

insert





union



Type of property	Number of properties	Bugs missed
Invariant	4	5
Postcondition	5	0
Metamorphic	16	0
Model-based	5	0

Effectiveness

```
prop FindPostPresent k v t =
  find k (insert k v t) === Just v
        May find
                          Will not find bugs
        bug in
                          in delete or
        find or
                          union
        insert
```

Effectiveness

2/3

prop_FindPostPresent k v t =
 find k (insert k v t) === Just v

insert



delete



union



Type of property	Number of properties	Bugs missed	Effectiveness
Invariant	4	5	38%
Postcondition	5	0	79%
Metamorphic	16	0	90%
Model-based	5	0	100%

```
=== prop_UnionPost from BSTSpec.hs:75 ===
Mean time to failure: 50.04595404595405
```

```
=== prop_InsertUnion fp
Mean time to failure
=== prop_DeleteUnion fr
Mean time to failure: 4
=== prop_UnionDeleteIns
Mean time to failure: 7
=== prop_UnionUnionAsso
Mean time to failure: 8
=== prop_FindUnion from
Mean time to failure
```

```
ec.hs:117 ===
        5374626
ogically equivalent!
      ec.hs:145 ===
      696303694
       BSTSpec.hs:167 ===
      32767233
      STSpec.hs:185 ===
      95104895
      .hs:206 ===
         72827
```

```
=== prop_UnionModel from pec.hs:290 ===
Mean time to failure: 8.368631368
```

```
prop_UnionPost t t' k =
  find k (union t t')
  ===
  (find k t <|> find k t')
```

```
prop UnionModel t t' =
  toList (union t t')
  List.sort
    (List.unionBy
       ((==) `on` fst)
       (toList t)
       (toList t'))
```

Mean time to failure

Property type	Min	Max	Mean
Postcondition	9.7	160	68
Metamorphic	1	401	61.6
Model-based	5	6.5	5.8

Averaged over seven bugs, and all properties of each type that detect the bugs



Model-based

- Easier to think of than postconditions
- Require fewer properties than metamorphic approach
- Are the most effective properties
- Find bugs fastest
- Complete specification

Metamorphic

- Do not require a model
- Easiest to write
- Good effectiveness